

# Performance and Reliability of High Density Flash EEPROMs Under CHISEL Programming Operation

S. Mahapatra<sup>1,3</sup>, S. Shukuri<sup>2</sup> and J. Bude<sup>3</sup>

<sup>1</sup>Department of Electrical Engineering, Indian Institute of Technology, Bombay 400076, India

<sup>2</sup>Semiconductor and Integrated circuits group, Hitachi Ltd., Tokyo, Japan

<sup>3</sup>Agere Systems, 600 Mountain Avenue, Murray Hill, NJ 07974, USA

Email: souvik@ee.iitb.ac.in

## Abstract

*We demonstrate CHISEL programming operation of fully scaled high-density flash EEPROMs. Single cell program and erase characteristics show reliable operation in terms of programming disturbs and cycling induced degradation. Program and erase operation of high-density arrays show a unique post-erase operation, tight threshold voltage distribution and over 10 years of data retention even after  $10^5$  program/erase cycles. Results are presented showing the feasibility of CHISEL programming operation for deeply scaled high-density flash EEPROMs.*

## 1. Introduction

CHannel Initiated Secondary Electron (CHISEL) injection (see Figure 1) was shown to be an excellent programming mechanism for flash EEPROMs [1-7]. It offers more efficient programming (faster speed, lower power) compared to standard Channel Hot Electron (CHE) operation. Initial results of CHISEL operation on large single cells ( $L_{FG} > 0.35\mu\text{m}$ ) and relatively small test arrays of 64k-bits show highly efficient self-convergent programming, a unique recovery procedure for over erased cells [8,9], a high margin for drain disturbs as well as good endurance for programmed and erased  $V_T$  up to  $10^5$  program/erase cycles, leading to reliable programming operation [4-7].

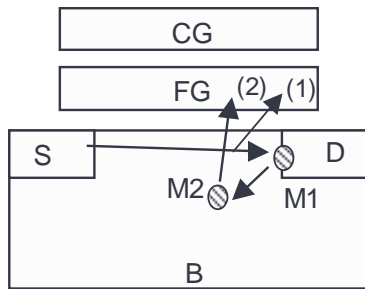


Figure 1. CHISEL injection mechanism. Channel electrons heated by lateral field undergo impact ionization M1. Holes generated from M1 flow to the substrate and undergo further impact ionization M2, which is enhanced under high transverse field (negative substrate bias). The secondary electrons (CHISEL) generated from M2 traverse towards the interface and get injected into the oxide (2) with greater energies than the lateral field heated channel electrons (1).

In this paper we demonstrate successful CHISEL programming operation of large arrays (up to 32M-bit) fabricated with fully scaled cells having  $L_{FG}=0.26\mu\text{m}$ . Single cell measurement results show program speed of  $1.3\mu\text{s}$  at  $V_D=4\text{V}$ ,  $V_T$  window closure of less than 1V and more than  $10^3$  program disturb margin after  $10^5$  program/erase cycles. Array measurement results show a unique over erase cell recovery procedure for post erase operation, tight program/erase  $V_T$  distribution and more than 10 years of data retention after  $10^5$  program/erase cycles. Results are presented that establish the viability of CHISEL programming operation of large arrays fabricated using scaled cells suitable for future high-density flash EEPROMs.

## 2. Device Fabrication

Fully scaled ( $W_{FG}=0.25\mu\text{m}$ )  $L_{FG}=0.26\mu\text{m}$  single cells and arrays (up to 32M-bit) have been fabricated using a  $0.18\mu\text{m}$  process involving state-of-the art modules like STI and self aligned source/drain contacts required in high-density memories. The cells have gate coupling of 0.6 and area of  $0.45\mu\text{m}^2$ . The SEM picture taken from a 32M-bit array is shown in Figure 2.

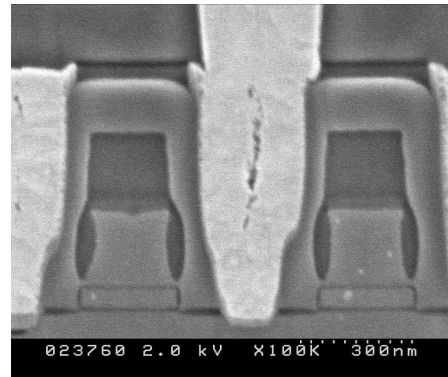


Figure 2. SEM of a 32M-bit array showing two adjacent cells. The floating (FG) and control (CG) gates and tunnel oxide (12nm) and inter poly dielectric (20nm) are shown.

## 3. Single cell operation

Program and erase transients of a  $L_{FG}=0.26\mu\text{m}$  cell before and after  $10^5$  program/erase cycles is shown in

Figure 3. The programming was performed at  $V_B = -2V$  with  $V_{CG}/V_D = 8/4V$  while channel erase was performed at  $V_{CG} = -22V$ . Threshold voltages at programmed and erased states were fixed at 5.4V and 1.8V respectively. The measured program ( $T_P$ ) and erase ( $T_E$ ) times are 1.3 $\mu s$  and 6.3ms (initial) and 2 $\mu s$  and 17ms (after  $10^5$  cycles) respectively. Compared to CHE programming at higher  $V_D$  for identical initial  $T_P$  (not shown in this paper), CHISEL shows much lower program time degradation with slightly higher erase time degradation due to cycling.

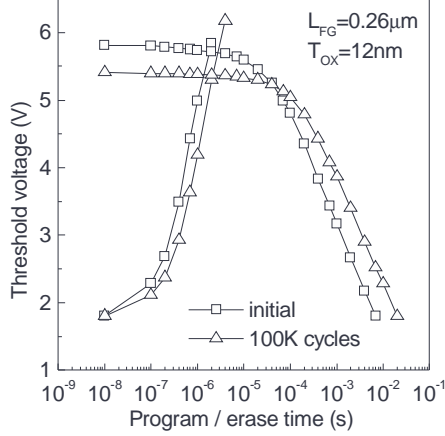


Figure 3. Program and erase transients before and after  $10^5$  repetitive program/erase cycles of a  $L_G = 0.26\mu m$  flash cell under CHISEL programming operation showing 1.3 $\mu s$  initial programming time at  $V_D = 4V$ .

Figure 4 shows  $V_T$  degradation in programmed ( $V_{TP}$ ) and erased ( $V_{TE}$ ) state (left y axis) and read current (right y axis) due to repetitive program/erase cycles. Programming pulse was applied for 1.3 $\mu s$  with  $V_{CG}/V_D/V_B = 8/4/-2V$ . Erase pulse was applied for 6.3ms with  $V_{CG} = -22V$ .  $V_T$  was defined at  $V_D/I_D = 0.8V/5\mu A$ , and read current was measured at erased state with  $V_{CG}/V_D = 4.5/0.8V$ .

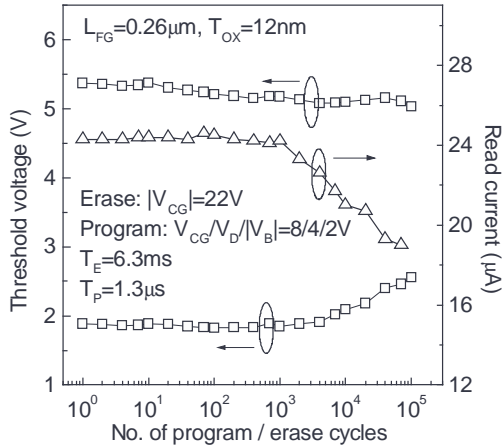


Figure 4. Cycling induced degradation in programmed and erased  $V_T$  (LHS) and read current (RHS) of a  $L_{FG} = 0.26\mu m$  cell. Programmed and erased  $V_T$  degradation is consistent with degradation in program and erase time as shown in Figure 3.

Note that program/erase cycling creates interface and oxide defects. Charges in these defects cause read current degradation due to scattering. They also act as a barrier to charge transfer and reduce program/erase current, causing degradation in  $V_{TP}$  (hence  $T_P$ ) and  $V_{TE}$  (and  $T_E$ ). In Figure 4, though some degradation can be seen for  $V_{TE}$  and read current,  $V_{TP}$  shows very little degradation. This is consistent with degradation in  $T_P$  and  $T_E$  (Figure 3) – relatively smaller degradation in  $T_P$  compared to  $T_E$ . Since CHISEL mechanism populates the high-energy tail of injected electrons they can easily overcome the increase in injection barrier, which explains the relatively lesser degradation in  $V_{TP}$  (and  $T_P$ ) for CHISEL operation [6]. This feature is unique to CHISEL programming operation that ensures very little overall window closure due to cycling.

Figure 5 shows charge gain and loss drain disturb measured before and after  $10^5$  program/erase cycles during CHISEL programming operation. Both cycling and disturb experiments were done at  $V_D = 4V$ . Charge gain and loss measurements were done respectively on erased and programmed cell while  $V_{CG}$  was held at 0V.

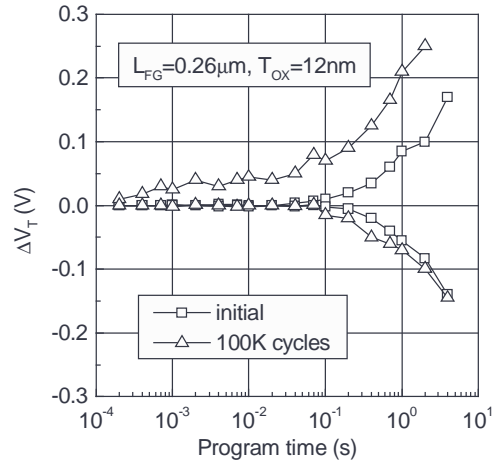


Figure 5. Charge gain and loss drain disturb for CHISEL programming operation of a  $L_{FG} = 0.26\mu m$  cell, before and after  $10^5$  cycles. After cycling, charge gain disturb becomes greater while there is no change for charge loss disturb. Programming condition same as Figure 3.

Charge gain disturb is due to electron injection into FG, is initiated by source-drain leakage and aided by uncharged FG (positive FG potential). Charge loss disturb is due to hole injection into FG, is initiated by band-band tunneling (BBT) and aided by charged FG (negative FG potential). After cycling, charge gain disturb is enhanced, while charge loss disturb remains largely unchanged. A maximum programming time of 2 $\mu s$  (after  $10^5$  cycles) and 128 cells per bit-line leads to total disturb time of 254 $\mu s$ . Figure 5 shows that more than  $10^3$  margin is available for both the disturb modes even after  $10^5$  cycles – clearly indicating that drain disturb is not a serious issue for CHISEL operation. When compared to CHE operation at identical  $T_P$  (not shown), CHISEL shows much lower charge gain disturb

with slightly higher charge loss disturb. This is attributed to lower source-drain leakage (higher  $V_T$  at  $V_B < 0$ ) but higher BBT (higher electric field across drain junction) for CHISEL operation.

#### 4. Array operation

Figure 6 shows the program, erase and post-erase  $V_T$  distribution measured on a 128K-bit block of a 4M-bit memory under CHISEL programming operation.

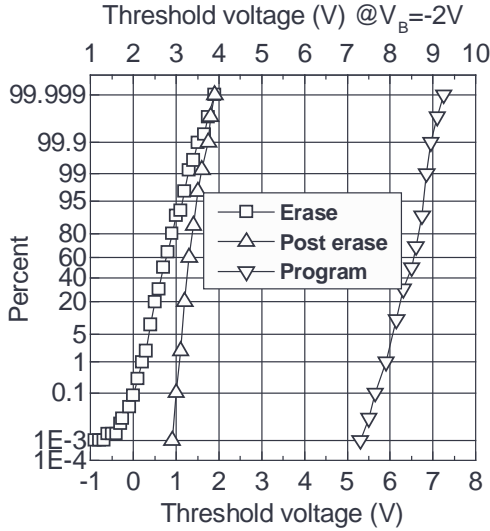


Figure 6. Program, erase and post erase  $V_T$  distribution of a 128k-bit block taken from a 4M-bit memory, made with fully scaled  $L_{FG}=0.26\mu m$  cells. Tight  $V_T$  distributions are obtained with CHISEL programming operation. Top x axis denotes the body effect induced  $V_T$  shift when  $V_B < 0$  is used for CHISEL programming.

Programming and post erase operations were done at  $V_D/V_B = 6/-2V$  with  $V_{CG}=8V$  and  $3V$  respectively. To achieve identical  $T_P$  as a single cell, array operation requires higher  $V_D$  to account for the voltage drop due to bit line series resistance. Channel erase was done at  $V_{CG}=-22V$ . A  $2\mu s$  programming pulse and  $6.3ms$  erase pulse were used. Post erase operation was performed using two  $10\mu s$  pulses with a single intermediate read verification scheme. The spread in  $V_T$  distribution is  $\sim 2V$  in the programmed state,  $\sim 3V$  after erase and only  $\sim 1V$  after post erase. The extremely tight post erase  $V_T$  distribution is obtained because of the self-converging nature of CHISEL programming which stops after the cell programs to a  $V_T$  as determined by  $V_{CG}=3V$ .

The unique recovery procedure of over-erased cells using CHISEL programming operation is described in Figure 7. Note from Figure 6 that  $V_T$  distribution after erase ("A", Figure 7) contains tail cells that are deeply over erased ( $V_T < 0$ ). These tail cells are difficult to recover under conventional CHE operation. However during CHISEL programming the application of a negative  $V_B$  cause  $V_T$  to increase via body effect shift (top x axis, Figure 6). The entire erased  $V_T$  distribution therefore shifts to  $V_T > 0$  ("B", Figure 7). CHISEL

programming can be performed to obtain the post-erase distribution ("C", Figure 7) which reverts to true post-erase distribution at  $V_B=0$  ("D", Figure 7).

Since CHISEL offers self convergent programming leading to tight  $V_T$  control and over erase is not an issue, cells can be programmed to relatively lower  $V_{TP}$  levels, opening up possibility for multi-level storage.

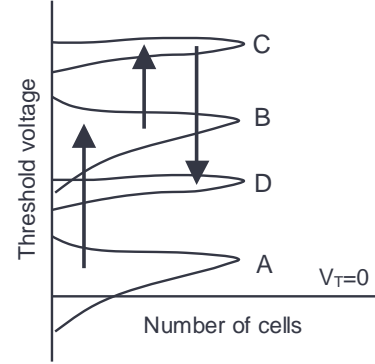


Figure 7. Over-erased cell recovery procedure under CHISEL programming operation. Body effect at  $V_B < 0$  shifts the  $V_T$  distribution from A to B. Post erase self convergent programming pushes the distribution to C, which reverts to D when  $V_B$  is set back to  $0V$ .

Figure 8 shows the programmed and erased  $V_T$  distribution during CHISEL programming operation on a 512k-bit block of a 32M-bit memory. Measurements were performed both before and after different amount of repetitive program/erase cycles, up to  $10^5$  cycles. Program, erase and post-erase conditions are identical to Figure 6, and the erased  $V_T$  distributions are shown after post-erase convergence.

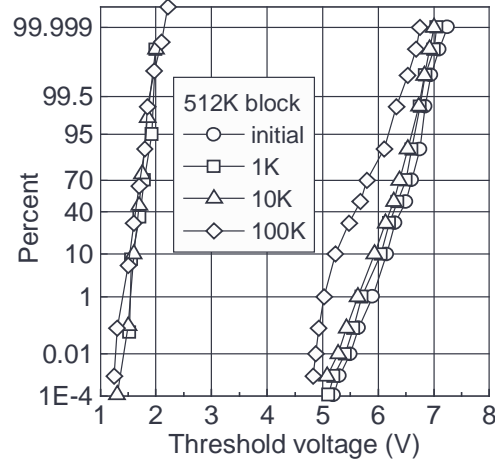


Figure 8. Cycling endurance of large array under CHISEL programming operation showing  $V_T$  distribution in programmed and erased state after different amount of repetitive program/erase cycles. Measurements were done on a 512k-bit block of a 32M-bit array. Extremely tight  $V_T$  distributions are obtained even after  $10^5$  cycles.

Note that even after  $10^5$  cycles the programmed and erased  $V_T$  distributions are tight, about  $\sim 2V$  and  $\sim 1V$  wide respectively. This proves that single cell results showing very little  $V_T$  window closure due to cycling for

CHISEL operation also holds true for large arrays. The extremely tight  $V_T$  distribution even after  $10^5$  cycles reinforces the possibility for multi-level storage.

## 5. Data retention

Data retention tests were performed on 32M-bit arrays after being subjected to  $10^5$  program/erase cycles. A positive substrate bias was applied during retention test to accelerate the charge loss from floating gate to substrate. Bit-by-bit measurements were performed (with grounded substrate) to obtain the  $V_T$  distribution after each stress interval. The worst bit of the programmed  $V_T$  distribution was identified and its  $V_T$  shift was extrapolated to estimate the worst bit among 1000 chips. Figure 9 shows the accelerated charge loss for the worst bit among 1000 chips of a 32M-bit array as a function of measurement time. For comparison, the non-accelerated charge loss data is also shown. The projected data loss (by linear extrapolation) ensures that the worst bit shows no bit failure in 10 years, even after the array has seen  $10^5$  CHISEL program/erase cycles.

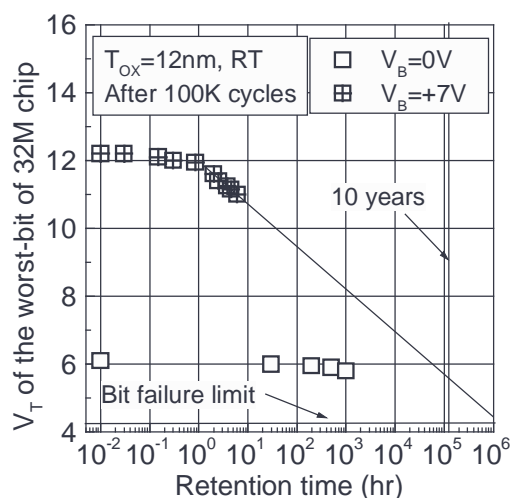


Figure 9. Estimated retention characteristics of the worst bit in 1000 32M-bit chips by using a 4M-bit test array operated in CHISEL mode. A positive substrate bias was applied to accelerate the charge loss from floating gate simulating programming to  $V_T=12V$ .

## 6. Conclusions

To summarize, we have demonstrated successful CHISEL programming operation on large arrays (up to 32M-bit) fabricated using fully scaled  $W_{FG}=0.25\mu m$ ,  $L_{FG}=0.26\mu m$  cells (area  $\sim 0.45\mu m^2$ ). The arrays were built using a state-of-the-art  $0.18\mu m$  technology having advanced modules like STI and self aligned source/drain contacts.

Single cell results show fast programming time of  $1.3\mu s$  ( $2\mu s$  after  $10^5$  cycles) at  $V_B=-2V$  and  $V_D=4V$ . The reliability of CHISEL operation in terms of

program/erase cycling induced degradation has been found to be extremely good.  $V_T$  window closure of less than 1V and more than  $10^3$  margin for programming drain disturb (for both charge gain and loss modes) have been achieved even after  $10^5$  cycles.

The unique over erase cell recovery procedure for CHISEL post erase operation (by employing the  $V_B<0$  induced body effect  $V_T$  shift) has been demonstrated using array measurements. We have shown that due to the self converging nature of CHISEL programming, an extremely tight  $V_T$  distribution of  $\sim 1V$  can be achieved by a post erase operation involving just 2 programming steps with a single intermediate read verification. Array cycling measurements also show extremely tight  $V_T$  distribution in program and erased state, even after  $10^5$  cycles. Finally, charge loss experiments on large arrays show more than 10 years data retention even after  $10^5$  cycles. Our results demonstrate the viability of CHISEL programming operation on large arrays suitable for future large density flash memories.

## References

- [1] J. D. Bude, "Gate current by impact ionization feedback in sub-micron MOSFET technologies", in *Proc. Symp. VLSI Technology*, p.101, 1995.
- [2] D. Esseni and L. Selmi, "A better understanding of substrate enhanced gate current in MOSFETs and flash cells – Part I: Phenomenological aspects", *IEEE Trans. Electron Devices*, vol.46, p.369, Feb.1999.
- [3] L. Selmi and D. Esseni, "A better understanding of substrate enhanced gate current in MOSFETs and flash cells – Part II: Physical analysis", *IEEE Trans. Electron Devices*, vol.46, p.376, Feb.1999.
- [4] J. D. Bude, A. Frommer, M. R. Pinto and G. R. Weber, "EEPROM/flash sub-3.0V drain-source bias hot carrier writing", in *IEDM Tech. Dig.*, p.989, 1995
- [5] J. D. Bude et. al., "Secondary electron flash – A high performance low power flash technology for  $0.35\mu m$  and below", in *IEDM Tech. Dig.*, p.279, 1997
- [6] J. D. Bude, M. R. Pinto and R. K. Smith, "Monte Carlo simulation of CHISEL flash memory cell", *IEEE Trans. Electron Devices*, vol.47, p.1873, Oct.2000.
- [7] M. Mastrapasqua, "Low voltage flash memory by use of a substrate bias", *Microelectron. Eng.*, vol.48, p.389, 1999.
- [8] C. Y. Hu et. al., "A convergence scheme for over erased flash EEPROMs using substrate enhanced hot electron injection", *IEEE Electron Device Lett.*, vol.11, p.500, 1995
- [9] C. Y. Hu et. al., "Substrate-current-induced hot electron (SCHE) injection: A new convergence scheme for flash memory", in *IEDM Tech. Dig.*, p.283, 1995.